A Patricia-Tree Approach for Frequent Closed Itemsets

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Abstract—In this paper, we propose an adaptation of the Patricia-Tree for sparse datasets to generate non redundant rule associations. Using this adaptation, we can generate frequent closed itemsets that are more compact than frequent itemsets used in Apriori approach. This adaptation has been experimented on a set of datasets benchmarks.

Keywords—Datamining, Frequent itemsets, Frequent closed itemsets, Sparse datasets.

I. INTRODUCTION

Efficient algorithms for mining frequent itemsets are crucial for mining association rules. Methods for mining frequent itemsets and for iceberg data cube computation have been implemented using a prefix-tree structure, namely FP-Tree, for storing compressed informations about frequent itemsets. As pointed out by Han [6] the FP-Tree loses his compactness on sparse datasets, but they still use this structure for mining frequent closed itemsets. In this paper we propose an enhanced version of Patricia-Tree structure that reduces considerably the size taken by an FP-Tree and the build time. This structure is especially suitable for sparse datasets.

Datamining algorithms based on frequent itemsets like Apriori suffer from two drawbacks: (i) multiple scans to a dataset to compute the frequency of itemsets; (ii) high number of generated association rules. To avoid these two drawbacks, many solutions are used, like formal analysis concepts, parallelism, data structures adapted to datamining and so on. In this paper, we propose an adaptation of the Patricia-Tree structure for sparse datasets to find frequent closed itemsets. Then, we experiment this new structure on different databases and compare it with the FP-Tree structure. The remaining of the paper is as follows: In section 2, we present briefly the main approaches used to generate association rules. Section 3 presents and compares the FP-Tree structure and the proposed adaptive Patricia-Tree for frequent closed itemsets. In section 4, we discuss some experimental results on dense and sparse datasets. Section 5 concludes the paper and gives some extensions of our work.

II. PREVIOUS WORK

A. Apriori-based Algorithms

The most generic frequent patterns mining algorithm is Apriori [2]. This algorithm is based on frequent itemsets that are generated from candidate itemsets [2]. Using this approach, a number of Apriori-based algorithms [1,3,4] have been developed. Among these algorithms, only those use a Hash-tree representation of the database are efficient [1].

B. Pattern Growth Algorithms

Han proposes a new technique for mining frequent itemsets without generating candidate itemsets [5]. It defines two new data structures: frequent pattern tree or FP-Tree to compact dense datasets and H-struct [6] to deal with sparse datasets solely. Later, Pietracaprina and Zandolin have proposed to use a better compressed tree, called Patricia-Tree [7].

C. Closed Itemsets Mining

To reduce the huge number of rules produced by algorithms based on frequent itemsets, Pasquier [8] proposes to generate only frequent closed (i.e. non redundant) patterns. The algorithms that generate frequent closed itemsets use either item-based data structures [8,9] or the FP-Tree structure [10,11].

III. PATRICIA-TREE VS FP-TREE

In this section, we compare the Patricia-Tree structure with FP-Tree in order to determine what is more accurate for different databases (dense or sparse).

A. FP-Tree

The FP-Tree structure consists of a set of prefix subtrees under a root node labeled as “null” and a header table containing frequent items [5]. Every header table entry points a node in the FP-Tree carrying the same item name and every node on the FP-Tree points to the next occurrence of this item.

B. Patricia-Tree

A Patricia-Tree is a compressed FP-Tree. We keep the same representation as an FP-Tree but we merge every parent node with his single child node when they have the same support value [7]. Contrarily to an FP-Tree node that represents a single item a Patricia-Tree node can represent several items.

C. Comparison

As pointed out by Han [6], the FP-Tree loses his compactness on sparse datasets, but they still use this structure
for mining frequent closed itemsets. The compactness of the FP-Tree is materialized by the merge of common prefixes for dense datasets. But for the sparse ones there is few prefixes shared by the transactions. In this case the number of FP-Tree nodes becomes close to the original dataset size (i.e. the sum of all transaction lengths).

For these reasons, we propose to adapt the Patricia-Tree proposed in [7] to generate frequent closed itemsets [13]. For example, consider the dataset represented by table 1 with

<table>
<thead>
<tr>
<th>TID</th>
<th>Items</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>A B D E F G H I</td>
</tr>
<tr>
<td>2</td>
<td>B C E L</td>
</tr>
<tr>
<td>3</td>
<td>A B D F H L</td>
</tr>
<tr>
<td>4</td>
<td>A B C D F G L</td>
</tr>
<tr>
<td>5</td>
<td>B G H L</td>
</tr>
<tr>
<td>6</td>
<td>A B D F I</td>
</tr>
</tbody>
</table>

Table 1: Sample Dataset D

*minsup* set to 3. Its representations by the FP-Tree and the Patricia-Tree are given respectively by figures 1 and 2.

[Fig. 1. FP-Tree for Dataset D]

[Fig. 2. Patricia-tree for Dataset D]

The above figures show that a Patricia-Tree is more compact than an FP-Tree. In fact, a dataset consisting of M transactions with aggregate size N can be represented through a Patricia-Tree of size at most equal to N + O(M) [7]. But for an FP-Tree when the dataset is highly sparse, the number of nodes may be close to N * M.

IV. EXPERIMENTAL RESULTS

In this section, we discuss results of several experimentations of our proposed data structure for different datasets. The goal of these experimentations is to find out the extent of different dataset properties that could affect the performance of Patricia-Tree and its relative performance with the FP-Tree [13].

Experiments were performed on a 500MHz Pentium PC with 320MB of memory, running on RedHat Linux 8.2.

Our version of Patricia-Tree was implemented in C and the FP-Tree was coded in C++ by Zhu [12].

For our experimentations, we have used several real and synthetic database benchmarks, publicly available at the FIMI\(^1\) workshop site. The PUMSB dataset contains census data. The MUSHROOM database contains characteristics of various species of mushrooms. The CONNECT dataset is derived from its game steps. The synthetic datasets T4011D100K and T2011D100K, obtained from IBM Almaden generator, mimic the transactions in a retailing environment.

Table 2 gives the characteristics of the real and synthetic datasets used in our evaluation. It shows the number of items, the average transaction length, the standard deviation of transaction lengths, and the number of records in each database.

The first experiment compares the FP-Tree and the Patricia-Tree build times for the different datasets.

As shown in Figures 3 and 5, these structures perform the same build time for real datasets. But Figure 4 shows that the build time of a Patricia-Tree is about two orders of magnitude faster than FP-Tree for synthetic datasets.

[Fig. 3. FP-Tree build time/Patricia-tree build time for real datasets.]

\(^1\)http://fimi.cs.helsinki.
The second experiment measures the memory space needed by FP-Tree and Patricia-Tree for the benchmark datasets. On real datasets Patricia-Tree consume more space than FP-Tree due to the over head taken by the number of items on a node and other data needed to tree traversal (Figures 6 and 8). For sparse datasets, Figure 7 shows that Patricia-Tree is more memory efficient.

V. CONCLUSION

We proposed an adaptation of the Patricia-Tree structure to find frequent closed itemsets. This adaptation allows to define more efficient datamining algorithms than those used FP-Tree structure. The experimentations of this adapted Patricia-Tree
structure have showed that it is more suitable for sparse datasets.
In the future, we plan to use this same structure to find sequential patterns and closed sequential patterns.

REFERENCES

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